A Compositional Semantics for OWL-S

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Introduction

- OWL is Web Ontology Language, proposed to W3C for Semantic Web
- OWL-S is 'service ontology', defining Semantic Web Services
- OWL-S process model describes formation of services by composition
- OWL-S process model therefore defines orchestration via workflow

Context

 OWL-S process model aims to capture common subset of workflow features

whereas

 WS-BPEL ends up with an all-inclusive superset of features

More Context

 YAWL attempts to capture all workflow 'patterns' in Petri net dialect

whereas

 Much process calculus-like work directly models specific features

'Composability' vs. Compositionality

- 'Composability' implies:
 - Existing (semantic) results should extend to new syntactic features

(not so for direct process model)

- Principle of compositionality:
 - Semantics of (syntactic) composite should derive from semantics of components (not so for Petri nets)

Aims

 Create a compositional model for OWL-S process model in general process algebra

but

Composable?

Compositional through what equivalence?

Approach

- Take existing model of coarsely (sequentially) interleaved dataflow (CONCUR03), which is one of OWL-S composite process types
- 'Compose' other OWL-S process types
- Apply existing notion of behavioural equivalence (temporal observation congruence)

(Generalised) OWL-S Processes

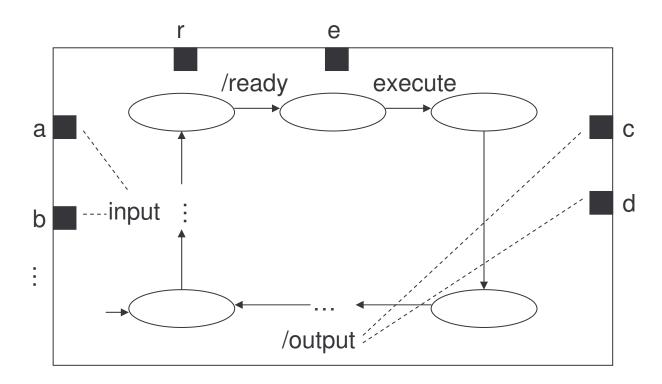
```
AtomicProcess ...
Process
         ::=
                CompositeProcess CProcess ...
                AnyOrder PerformanceList |
CProcess ::=
                Sequence PerformanceList
                Split PerformanceList
                SplitJoin PerformanceList |
                ChooseOne PerformanceList |
                IfThenElse Performance
                            Performance |
                RepeatWhile Peformance |
                RepeatUntil Performance
PerformanceList Performance
                PerformanceList; Performance |
                PerformanceList; Connect ...
Performance ::= Perform Process
```

Existing Model

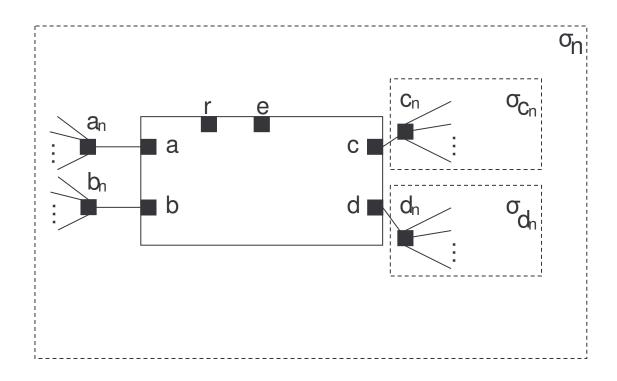
- Take automata describing interfaces of components
- Compose agent representing participation in *global synchronisation* to form instance
- Compose instances together, in model aware of communication-style (local) and global synchronisations
- Compare for conformance to interface (automaton) assigned to composite

'Interface Automata'

Generally:



Instantiation



Basis (Regular CCS) ...

$$\mathcal{E} ::= \mathbf{0} \mid \alpha . \mathcal{E} \mid \mathcal{E} + \mathcal{E} \qquad | \mu X . \mathcal{E} \mid X$$

$$a, \overline{a}, b, \overline{b}, \dots \in \Lambda \cup \overline{\Lambda}$$

$$\alpha, \beta, \dots \in \Lambda \cup \Lambda \cup \{\tau\}$$

Act
$$\frac{-}{\alpha.P\overset{\alpha}{\to}P}$$
 Sum1
$$\frac{P\overset{\alpha}{\to}P'}{P+Q\overset{\alpha}{\to}P'}$$

Sum2
$$\frac{Q \stackrel{\alpha}{\rightarrow} Q'}{P + Q \stackrel{\alpha}{\rightarrow} Q'}$$

$$\operatorname{Rec} \ \frac{E\{\mu X.E/X\} \stackrel{\alpha}{\to} P}{\mu X.E \stackrel{\alpha}{\to} P}$$

Basis (CCS) ...

$$\mathcal{E} ::= \mathbf{0} \mid \alpha.\mathcal{E} \mid \mathcal{E} + \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \qquad \mid \mu X.\mathcal{E} \mid X$$
$$a, \overline{a}, b, \overline{b}, \dots \in \Lambda \cup \overline{\Lambda}$$
$$\alpha, \beta, \dots \in \Lambda \cup \Lambda \cup \{\tau\}$$

$$\operatorname{Rec} \ \frac{E\{\mu X.E/X\} \stackrel{\alpha}{\to} P}{\mu X.E \stackrel{\alpha}{\to} P}$$

+ Deterministic Time ...

$$\mathcal{E} ::= \mathbf{0} \mid \alpha.\mathcal{E} \mid \mathcal{E} + \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \mid \sigma(\mathcal{E}) \mid \mu X.\mathcal{E} \mid X$$

$$a, \overline{a}, b, \overline{b}, \dots \in \Lambda \cup \overline{\Lambda}$$

$$\alpha, \beta, \dots \in \Lambda \cup \Lambda \cup \{\tau\}$$

$$\gamma, \delta \dots \in \Lambda \cup \overline{\Lambda} \cup \{\tau, \sigma\}$$

+ Maximal Progress (≈TPL) ...

$$\mathcal{E} ::= \mathbf{0} \mid \qquad \alpha.\mathcal{E} \mid \mathcal{E} + \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \rfloor \sigma(\mathcal{E}) \mid \mu X.\mathcal{E} \mid X$$

$$a, \pi, b, \overline{b}, \dots \in \Lambda \cup \overline{\Lambda}$$

$$\alpha, \beta, \dots \in \Lambda \cup \Lambda \cup \{\tau\}$$

$$\gamma, \delta \dots \in \Lambda \cup \overline{\Lambda} \cup \{\tau, \sigma\}$$

Act
$$\frac{-}{\alpha.P\overset{\alpha}{\rightarrow}P}$$
 TO2 $\frac{P\overset{\alpha}{\rightarrow}P'}{[P]\sigma(Q)\overset{\alpha}{\rightarrow}Q}$ TO2 $\frac{P\overset{\alpha}{\rightarrow}P'}{[P]\sigma(Q)\overset{\alpha}{\rightarrow}P'}$ Sum1 $\frac{P\overset{\alpha}{\rightarrow}P'}{P+Q\overset{\alpha}{\rightarrow}P'}$ Sum2 $\frac{Q\overset{\alpha}{\rightarrow}Q'}{P+Q\overset{\alpha}{\rightarrow}Q'}$ Sum3 $\frac{P\overset{\alpha}{\rightarrow}P'}{P+Q\overset{\sigma}{\rightarrow}P'+Q'}$ Com1 $\frac{P\overset{\alpha}{\rightarrow}P'}{P|Q\overset{\alpha}{\rightarrow}P'|Q}$ Com2 $\frac{Q\overset{\alpha}{\rightarrow}Q'}{P|Q\overset{\alpha}{\rightarrow}P|Q'}$ Com3 $\frac{P\overset{\alpha}{\rightarrow}P'}{P|Q\overset{\pi}{\rightarrow}P'|Q'}$ Com4 $\frac{P\overset{\alpha}{\rightarrow}P'}{P|Q\overset{\sigma}{\rightarrow}P'|Q'}$ Patient $\frac{-}{a.P\overset{\sigma}{\rightarrow}P}$ Idle $\frac{-}{\mathbf{0}\overset{\sigma}{\rightarrow}\mathbf{0}}$ Rec' $\frac{E\{\mu X.E/X\}\overset{\gamma}{\rightarrow}P}{P}$

+ 'Stalling'(=PMC's 0)...

$$\mathcal{E} ::= \mathbf{0} \mid \Delta \mid \alpha.\mathcal{E} \mid \mathcal{E} + \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \mid \sigma(\mathcal{E}) \mid \mu X.\mathcal{E} \mid X$$



N.B. Equivalent to adding insistent communication, since:

$$\underline{a} \cdot P \stackrel{\text{def}}{=} a \cdot P + \Delta$$

$$a, \overline{a}, b, \overline{b}, \dots \in \Lambda \cup \overline{\Lambda}$$

 $\alpha, \beta, \dots \in \Lambda \cup \Lambda \cup \{\tau\}$

$$\gamma, \delta \cdots \in \Lambda \cup \overline{\Lambda} \cup \{\tau, \sigma\}$$

+ Multiple Clocks (a la PMC, CSA)...

$$\mathcal{E} ::= \mathbf{0} \mid \Delta \mid \alpha.\mathcal{E} \mid \mathcal{E} + \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \mid \sigma(\mathcal{E}) \mid \mu X.\mathcal{E} \mid X$$

$$a, \overline{a}, b, \overline{b}, \dots \in \Lambda \cup \overline{\Lambda}$$

$$\alpha, \beta, \dots \in \Lambda \cup \Lambda \cup \{\tau\}$$

$$\rho, \sigma, \dots \in \mathcal{T}$$

$$\gamma, \delta \dots \in \Lambda \cup \overline{\Lambda} \cup \{\tau\} \cup \mathcal{T}$$

+ Hiding (= CaSE)

$$\mathcal{E} ::= \mathbf{0} \mid \Delta \mid \alpha.\mathcal{E} \mid \mathcal{E} + \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \rfloor \sigma(\mathcal{E}) \mid \mu X.\mathcal{E} \mid X$$

$$a, \overline{a}, b, \overline{b}, \dots \in \Lambda \cup \overline{\Lambda}$$

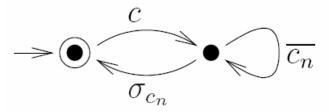
$$\alpha, \beta, \dots \in \Lambda \cup \Lambda \cup \{\tau\}$$

$$\rho, \sigma, \dots \in \mathcal{T}$$

$$\gamma, \delta \dots \in \Lambda \cup \overline{\Lambda} \cup \{\tau\} \cup \mathcal{T}$$

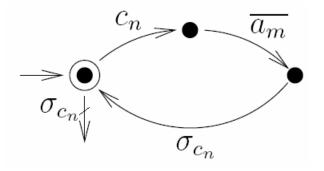
Compositional Broadcast

Broadcast



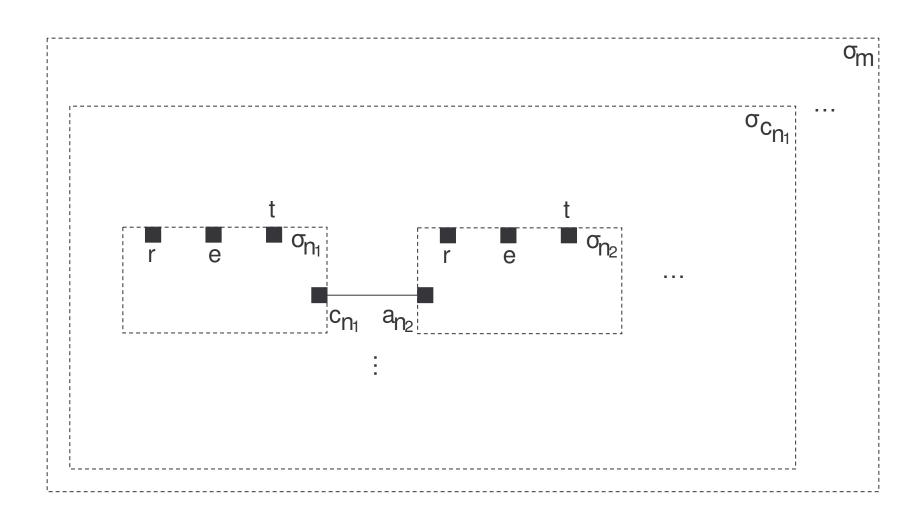
$$\mu X.\underline{c}_{\sigma_{c_n}}.\ \mu Y.\lfloor \underline{\overline{c_n}}.Y \rfloor \sigma_{c_n}(X)$$

Connection



$$\mu X.c_n.\overline{\underline{a_m}}.\underline{\sigma_{c_n}}.X$$

Composition



Conclusion

- Modulo small extension, calculus allows 'composition' of OWL-S process types
- Theoretical results:
 - Temporal observation congruence holds
 - Full abstraction holds
 - To do:
 - extension of algebraic theory
- Practical results
 - Implementation in Haskell
 - To do:
 - extend partition refinement
 - implementation in LISP

Further Work

- Fix cashew-s as a rich language for choreography (WSMO insists service interface provide both orch & chor)
- Establish expressiveness of cashew-nuts to give semantics to orchestration and choreography
- Investigate temporal observation congruence as a conformance test between orch & chor
- Now part of DIP (European Integrated) Project, therefore providing input to WSMO