

Towards a reconciliation of ChOrch in IRS, Configurator and WSMO

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Input

- Considers following work:
 - IRS
 - Choreography engine based on ASMs
 - Orchestration engine based on partial OWL-S

Configurator

- Partial UML2AD workflow as choreography
- UML2AD composes these to make orchestrations

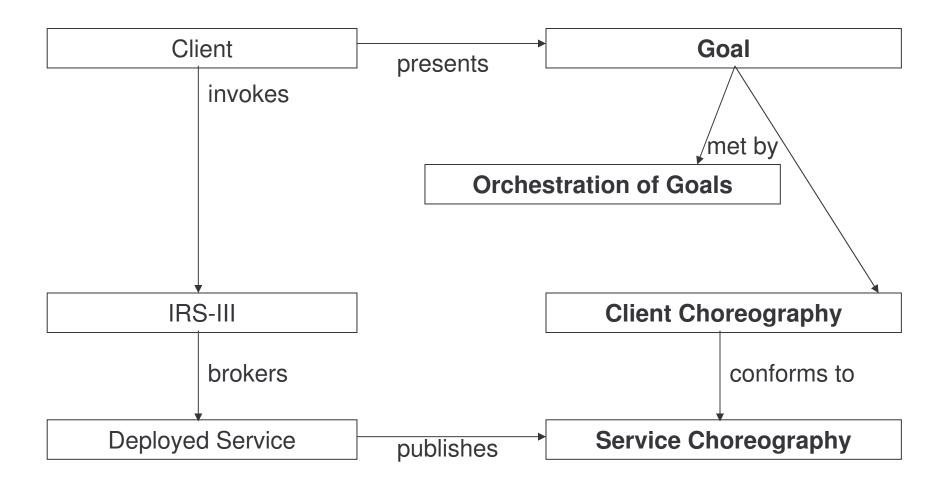
Cashew

- IO automata as choreographies
- OWL-S as orchestration (visualised in UML2AD), composes these, engine in Haskell via process algebraic semantics...

- WSMO

ASMs are everything (somehow...)

IRS View of ChOrch



Views on IRS

- Client choreography (currently an ASM) can be viewed as partial workflow
- Composition by orchestration of goals can be compared to workflow composition
- Answers Cashew criticism that
 - OWL-S
 - tackles only 'operation composition'
 (since it combines operations to make 'scripts' over a service that are atomic workflow tasks)
 - ignores challenges of (service) choreography (since it encapsulates dependencies between operations)
 - Configurator
 - considers only one (client's intention) interaction (claim: partial workflows can be viewed as client choreographies)

Requirements

- Need to achieve:
 - Reconciliation of viewpoints (IRS reference implementation, Configurator work, WSMO) on paper;
 - Demonstrator that convincingly executes an example illustrating this, i.e.
 - import from Configurator to IRS-III
 - orchestration engine in IRS-III (and interface with choreography engine)
 - export as ASMs
 - orchestration engine in WSMX

Proposal

- Build an ontology fragment:
 - representing workflow patterns;
 - structured as per OWL-S (process model);
 - adapted to capture UML idioms.
- Represent in this language:
 - orchestration of goals (goals as tasks);
 - client choreographies (operations of deployed service as tasks - restricted fragment?).
- Translate via Cashew to (control state) ASMs via process algebra...

Cashew Process Algebra Syntax and Semantics

$$\mathcal{E} ::= \mathbf{0} \mid \Delta \mid \alpha.\mathcal{E} \mid \mathcal{E} + \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \mid \mathcal{E} \mid \sigma(\mathcal{E}) \mid \mu X.\mathcal{E} \mid X$$

$$a, \overline{a}, b, \overline{b}, \dots \in \Lambda \cup \overline{\Lambda}$$

$$\alpha, \beta, \dots \in \Lambda \cup \Lambda \cup \{\tau\}$$

$$\rho, \sigma, \dots \in \mathcal{T}$$

$$\gamma, \delta \dots \in \Lambda \cup \overline{\Lambda} \cup \{\tau\} \cup \mathcal{T}$$